EECS2011 Fundamentals of Data Structures (Winter 2022)

Q&A - Lectures 7

Wednesday, March 9

Announcements

- Lecture W8 released
  - + General Trees vs. Binary Trees
  - + Implementing a Generic, General Tree in Java
  - + Terminology and Mathematical Properties
- Assignment 2 released
- Assignment 1 results released
- Written Test 2 coming soon

will be covered in WT2 → more challenging than WT1

Part B1-B4.

task 2: O(n)

1. office hours
2. submit regrading request

If we made a github repository of a past course, like eecs1022, is it ok if we share some projects that we did with employers or is it a breach for academic integrity.

EECS2011
↳ private before official end of the course (end of april / early may)

On week 6 slides (ADTs-stack....),
why do we declare our methods as static on pages 28 and 30 ?

```
public static <E> void reverse(E[] a) {
  Stack<E> buffer = new ArrayStack<E>();
  for (int i = 0; i < a.length; i ++) {
    buffer.push(a[i]);
  }
  for (int i = 0; i < a.length; i ++) {
    a[i] = buffer.pop();
  }
}
```

1. method does not need to use any attributes in StackUtilities class (this method only operates on the parameter).

2. StackUtilities.reverse(...);

StackUtilities su = new ...;
su.reverse(...);
↳ warning!

× ArrayList

ListOrderEPs ×

public MyClass < ___ , ___ , ... >

what names can we use as generic parameters.

any name that does not clash with:

(1) library class names

(2) classes created under the same project

E, I, K, V (single character in upper case)

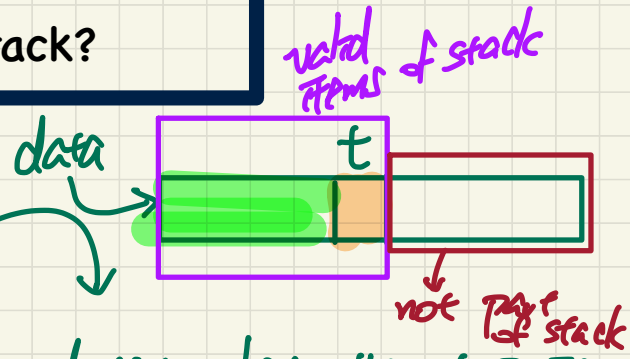
G, H, I, J, ...

Why do we need to assign null to the returned array element when we pop from an array stack?
Can't we just decrement the head of the stack?

```
public class ArrayStack<E> implements Stack<E> {
  private final int MAX_CAPACITY = 1000;
  private E[] data;
  private t; /* index of top */
  public ArrayStack() {
    data = (E[]) new Object[MAX_CAPACITY];
    t = -1;
  }

  public int size() { return (t + 1); }
  public boolean isEmpty() { return (t == -1); }

  public E top() {
    if (isEmpty()) { /* Precondition Violated */ }
    else { return data[t]; }
  }
  public void push(E e) {
    if (size() == MAX_CAPACITY) { /* Precondition Violated */ }
    else { t ++; data[t] = e; }
  }
  public E pop() {
    E result;
    if (isEmpty()) { /* Precondition Violated */ }
    else { result = data[t]; data[t] = null; t --; }
    return result;
  }
}
```



1. t indicates where the top is

2. items stored at indices > t are not valid stack items
   ↳ they should be detached ASAP whenever possible

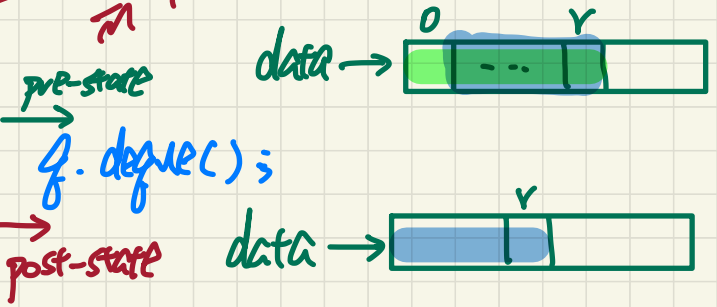
my question is about the dequeue method in circular array. The dequeue method in in circular array has to be void rather than returning a new array like arrayqueue?

```
public class ArrayQueue<E> implements Queue<E> {
  private final int MAX_CAPACITY = 1000;
  private E[] data;
  private int r; /* rear index */
  public ArrayQueue() {
    data = (E[]) new Object[MAX_CAPACITY];
    r = -1;
  }
  public int size() { return (r + 1); }
  public boolean isEmpty() { return (r == -1); }
  public E first() {
    if (isEmpty()) { /* Precondition Violated */ }
    else { return data[0]; }
  }
  public void enqueue(E e) {
    if (size() == MAX_CAPACITY) { /* Precondition Violated */ }
    else { r ++; data[r] = e; }
  }
  public E dequeue() {
    if (isEmpty()) { /* Precondition Violated */ }
    else {
      E result = data[0];
      for (int i = 0; i < r; i ++) { data[i] = data[i + 1]; }
      data[r] = null; r --;
      return result;
    }
  }
}
```

Q1. Why do we need C.A.?

Q2. Why do we need D.A?

this header was declared in the Queue interface



Hello professor. Can you please elaborate a bit on how to implement a queue using two stacks? Thanks.

```
public class StackQueue<E> implements Queue<E> {
  private Stack<E> inStack;
  private Stack<E> outStack;
  ...
}
```

static type

- For `size`, add up sizes of `inStack` and `outStack`.
- For `isEmpty`, are `inStack` and `outStack` both empty?
- For `enqueue`, ***push*** to `inStack`.
- For `dequeue`:
  - ***pop*** from `outStack`
    If `outStack` is empty, we need to first ***pop*** <u>all</u> items from `inStack` and ***push*** them to `outStack`.

✓

push to inStack

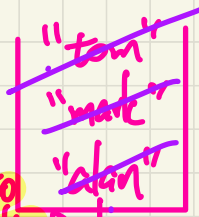
```
Queue<String> q = new StackQueue<>();
q.enqueue("alan");
q.enqueue("mark");
q.enqueue("tom");
String first = q.dequeue();
String second = q.dequeue();
String queue = q.dequeue();
```

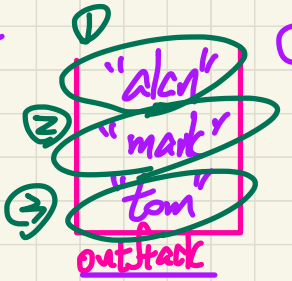
① first pop all items in inStack then push to outStack

client → FIFO

front

Supplier → Simulate FIFO using two LIFO stacks: inStack





② pop off outStack

One pattern:

enqueue ... enqueue dequeue ... dequeue

n operations

$O(1)$ $O((n-i)\cdot 1) = O(n)$ Amortized RT: $O\left(\frac{n+(n-1)}{n}\right) = O(1)$

Amortized RT of StackQueue

dequeue

worst case:

$O(n)$: pushing to outStack items popped from inStack

However, such transfer of items from in- to out-stack happens rarely. ⇒ ART: $O(1)$

Amortized (doubling) RT of dynamic arrays (for queue)

enqueue

worst case:

$O(n)$: copying items to the resized array

However, such resizing of array is not needed often

⇒ A.RT: $O(1)$

It seems that we can implement this without recursion as well.
We can simply keep updating the "from" and "to" in a loop.
Why do we not do that?

```
boolean binarySearch(int[] sorted, int key) {
  return binarySearchH(sorted, 0, sorted.length - 1, key);
}
boolean binarySearchH(int[] sorted, int from, int to, int key) {
  if (from > to) { /* base case 1: empty range */
    return false; }
  else if(from == to) { /* base case 2: range of one element */
    return sorted[from] == key; }
  else {
    int middle = (from + to) / 2;
    int middleValue = sorted[middle];
    if(key < middleValue) {
      return binarySearchH(sorted, from, middle - 1, key);
    }
    else if (key > middleValue) {
      return binarySearchH(sorted, middle + 1, to, key);
    }
    else  { return true; }
  }
}
```

Exercise: implement binary search using a loop

binary search tree